

W1001
28/12/14

Prob. Methods in Comb.

Azuma's ineq.: Let x_0, \dots, x_m be a martingale with $|x_{i+1} - x_i| \leq 1$ for each $i < m$. Then for every $x \geq 0$, $\mathbb{P}(x_m > \mathbb{E}x_m + x\sqrt{m}) \leq e^{-x^2/2}$

Chromatic Number of $G(n, 1/2)$

We know $\chi(G) \geq \frac{|V(G)|}{\alpha(G)}$ so we'd first like to compute $\alpha(G)$ for $G \sim G(n, 1/2)$. Let $f(k) = \mathbb{E}[\# \text{copies of } K_k \text{ in } G]$ (equivalent to trying to find an empty subgraph).

Then $f(k) = \binom{n}{k} 2^{-\binom{k}{2}}$ so $\frac{f(k-1)}{f(k)} = \frac{n-k}{k-1} 2^{-k}$. Roughly

Let $k_0 = \min\{k \geq 1 \mid f(k) < 1\}$. Claim - a. a. s. $w(G) \in \{k_0 - 1, k_0, k_0 + 1\}$

1) $k_0 \leq (2+o(1)) \log_2 n$ since $f(k) \leq \left(\frac{en}{k}\right)^k 2^{-\frac{k-1}{2}k} \ll 1$ for $2^{\frac{k-1}{2}} > \frac{en}{k}$, which happens roughly at $2 \log_2 n$.

2) $k_0 \geq (2-o(1)) \log_2 n$ since $f(k) \geq \left(\frac{n-k}{k}\right)^k 2^{-\frac{k-1}{2}k} \gg 1$ for $\frac{n-k}{k} > 2^{\frac{k-1}{2}} \Leftrightarrow k = (2-\epsilon) \log_2 n$ (for large enough n - $\forall \epsilon > 0 \exists N, \forall n > N \dots$).

So, $\frac{f(k_0+1)}{f(k_0)} \ll n^{-1+o(1)}$, so $f(k_0 \pm 1) \leq n^{-1+o(1)}$ and by the

first moment method G doesn't contain $K_{k_0 \pm 1}$ a. a. s.,

so $w(G) \leq k_0$ a. a. s. (equivalently $\alpha(G) \leq k_0$). We conclude

$\chi(G) \geq (1+o(1)) \frac{n}{2 \log_2 n}$. We prove (Bollobás) that in fact equality holds above.

Lemma $\mathbb{P}(w(G(n, 1/2)) < k_0 - 4) \leq e^{-n^{2+o(1)}}$. Equality above follows: $m = \lfloor \frac{n}{2 \log_2 n} \rfloor$. Suffices to prove a. a. s. $|S| = m \Rightarrow \alpha(G[S]) \geq (2-o(1)) \log_2 n$.

by lemma & union bound. $G \sim G(n, 1/2)$ so by the lemma $\mathbb{P}(\alpha(G[S]) < k_0(m) - 4) \leq e^{-m^{2+o(1)}} = e^{-n^{2+o(1)}}$.

but # sets $S \leq 2^n$, so the prob. the above holds for any S (by union bound) $\geq 1 - 2^n e^{-n^{2+o(1)}} \geq 1+o(1)$.

pf. of Lemma $\chi = \text{Max. size of a family of edge-disjoint } K_k \text{ cliques } (k = k_0 - 1)$ claim $\mathbb{E}(\chi) \geq (1+o(1)) \frac{n}{2}$

pf. claim \Rightarrow lemma Let $Y_0, \dots, Y_{\binom{n}{2}}$ be the edge exposure martingale, Y is Lipschitz so by Azuma

$$P(W(G) \leq k_1) = P(Y_{\binom{n}{2}} = 0 = EY - EY) \leq e^{-\frac{(EY)^2}{2\binom{n}{2}}} = e^{-\frac{n^2}{2 \log^2 n}}$$

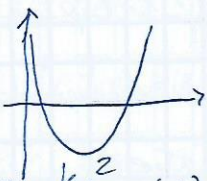
Now, $f(k_2) \geq f(k_0 - 5) = f(k_0 - 1) \prod_{i=1}^4 \frac{f(k_0 - i - 1)}{f(k_0 - i)} = f(k_0 - 1) \cdot n^{4+o(1)} \geq n^{4+o(1)}$

Define $W = \#$ intersecting pairs of cliques.

Fix $0 < q < 1$ and take C a collection of cliques obtained by selecting each K_{k_2} in G with prob. q . C' is given by deleting a clique from each intersecting pair in C .

$E[C'] \geq E[C] - q^2 E[W] = q f(k_2) - q \frac{E[W]}{2}$. Now,

$E[W] = f(k_2) \sum_{s=2}^{k_2-1} \binom{k_2}{s} \binom{n-k_2}{k_2-s} 2^{-\binom{k_2}{2} - \binom{s}{2}}$. Then $g(0) = (1-o(1)) f(k_2)^2$ and

$\frac{g(s+1)}{g(s)} = \frac{k_2 - s}{s+1} \cdot \frac{k_2 - s}{n - 2k_2 + s + 1} 2^s \rightarrow$  So this is

$f(k_2) (1+o(1)) (g(2) - g(k_2 - 1))$. $g(2) = \frac{k_2^2}{n} g(1) = \frac{k_2^4}{n^2} g(0) = f(k_2) \frac{k_2^4}{n^2} (1+o(1))$

and $g(k_2 - 1) \leq k_2 n 2^{-k_2 + 1} \leq \frac{k_2}{n} < \frac{k_2^4}{n^2} f(k_2)$. So,

$E[W] = (1+o(1)) \frac{k_2^4}{n^2} f(k_2)^2$. Pick $q = \frac{n^2}{k_2^4 f(k_2)} < 1$, so we get

$E[C'] = \frac{n^2}{k_2^4 f(k_2)} (f(k_2) - \frac{n^2}{k_2^4 f(k_2)} \frac{1}{2} \frac{k_2^4}{n^2} f(k_2)^2) = \frac{n^2}{2k_2^4}$ as promised.

Thm. Let $p = n^{-\alpha}$, $\alpha \leq 5/6$. Then $\exists u = u(n, p)$ s.t. a.a.s.

$u \leq \chi(G(n, p)) \leq u+3$.

Lemma Fix $\hat{\alpha}, \hat{c}$, $p = n^{-\hat{\alpha}}$. Then a.a.s. every subgraph of $G(n, p)$ with $\leq c\sqrt{n}$ vxs is 3-colorable.

pf. Enough that $\forall |S| \leq c\sqrt{n} = 2e_G(S) < \frac{3|S|}{2}$. $P(\text{this won't hold for some } S) \leq \sum_{s=4}^{c\sqrt{n}} \binom{n}{s} \binom{\binom{s}{2}}{\lfloor s/2 \rfloor} p^{\frac{3s}{2}}$. $b(s) \leq \left(\frac{en}{s}\right)^s \left(\frac{es^2 p}{3s}\right)^{\frac{3s}{2}} = \left(\frac{e^{5/2} n^5 p^{3/2}}{3^{3/2}}\right)^s \leq (cn^{-\delta})^s$ for some $\delta(\alpha)$ (since $\frac{5}{4} - \frac{3}{2} \frac{5}{6} = 0$). smallest

pf of thm. Let $\epsilon > 0$ be arbitrarily small and let u be such that

$P(\chi(G) \leq u) \geq \epsilon$. $Y = \min\{|S| \mid G-S \text{ is } u\text{-colorable}\}$.

Let Y_0, \dots, Y_n be vxs. exposure martingale. It's vertex-Lipschitz to by Azuma, $P(Y - EY \leq -\lambda\sqrt{n}) \leq e^{-\lambda^2/2}$. Fix λ

such that $e^{-\lambda^2/2} \leq \epsilon$. Then $P(Y \leq EY - \lambda\sqrt{n}) \leq \epsilon \leq P(Y/C) \leq u$.

$\mathbb{E}Y < \lambda_0 \sqrt{n}$, claim with prob $\geq 1 - 3\epsilon$, $\chi(G) \in \{u, u+1, u+2, u+3\}$.

$\mathbb{P}(\chi(G) < u) < \epsilon$, $\mathbb{P}(Y > \lambda_0 \sqrt{n}) < \epsilon$ so by lemma,

$\mathbb{P}(\text{Some set isn't 3-col.}) < \epsilon$, so we can color the rest using ≤ 3 cols. - the total probability that anything fails $\leq 3\epsilon$.

Ex. B arbitrary normed space, $v_1, \dots, v_n \in B \forall i$ $\|v_i\| \leq 1$. Let $\epsilon_1, \dots, \epsilon_n$ i.i.d. $\mathbb{P}(\epsilon_i = \pm 1) = 1/2$. $X = \left\| \sum_{i=1}^n \epsilon_i v_i \right\|$. $\mathbb{P}(X - \mathbb{E}X > \lambda \sqrt{n}) < e^{-\lambda^2/2}$.

pf. Note that (by tri. ineq.) changing $\epsilon_i \rightarrow -\epsilon_i$ changes X by at most 2. Therefore define $x_i = \mathbb{E}[X | \epsilon_1, \dots, \epsilon_i]$ so that x_0, \dots, x_n form a martingale. Therefore $|x_{i+1} - x_i| \leq 2$. But for the bound we need 1-Lipschitz! We may get a better estimate on $|x_{i+1} - x_i|$: Fix $\epsilon \in \{-1, 1\}^n$ and let ϵ' be ϵ with ϵ_{i+1} negated. Then $x_i(\epsilon) = \frac{1}{2}(x_{i+1}(\epsilon) + x_{i+1}(\epsilon'))$, so $|x_i(\epsilon) - x_{i+1}(\epsilon)| = |x_{i+1}(\epsilon) - x_{i+1}(\epsilon')| \leq 1$.

Indep. Bounded Differences ineq. (no proof, McDiarmid's "Concentration") Let x be a vector of indep. r.v. Suppose each $x_i \in A_i$ and $f: \prod_{i=1}^n A_i \rightarrow \mathbb{R}$, such that $|f(x) - f(x')| \leq C_i$ if x, x' differ only on the i^{th} coordinate. Then, letting $\mu = \mathbb{E}[f(x)]$, for every $t \geq 0$, $\mathbb{P}(f(x) - \mu \geq t) \leq e^{-\frac{2t^2}{\sum C_i^2}}$.

Applications: Isoperimetric ineq. of the hypercube

$\Omega = \Omega_1 \times \dots \times \Omega_n$. Given $x, y \in \Omega$ the Hamming dist. $d_H(x, y) = \#\{i \mid x_i \neq y_i\}$.

Thm. Let (x_1, \dots, x_n) be ind. r.v., $x_i \in \Omega_i$, $A \subseteq \Omega$ arbitrary. Then for any $t \geq 0$, $\mathbb{P}(x \in A) \mathbb{P}(d_H(x, A) \geq t) \leq e^{-t^2/2n}$.

Cor. $A \subseteq \{0, 1\}^n$, $|A| \geq \epsilon \cdot 2^n$, $\frac{|A_t|}{|A|} \geq 1 - \frac{e^{-t^2/2n}}{\epsilon}$.

pf. Let $p = \mathbb{P}(x \in A)$, $\mu = \mathbb{E}[d_H(x, A)]$. By IBDI for $t \geq 0$, $\mathbb{P}(d_H(x, A) - \mu \geq t) \leq e^{-2t^2/n}$. Note $p = \mathbb{P}(d_H(x, A) - \mu \leq -\mu) \leq e^{-2\mu^2/n}$, so $\mu \leq \sqrt{\frac{n}{2} \log \frac{1}{p}} = t_0$. Then

$\mathbb{P}(d_H(x, A) \geq t_0 + t) \leq e^{-t^2/n}$. For $t \geq t_0$, $\mathbb{P}(d_H(x, A) \geq t) \leq e^{-2(t-t_0)/n}$.
 If $t \geq 2t_0$, $(t-t_0)^2 \geq t^2/4$ so $\mathbb{P}(d_H(x, A) \geq t) \leq e^{-t^2/2n}$. If $t < 2t_0$,
 recall $p = \mathbb{P}(x \in A) = e^{-2t_0^2/n}$. Therefore $\min\{\mathbb{P}(x \in A), \mathbb{P}(d_H(x, A) \geq t)\}$
 is $\leq e^{-t^2/2n}$.

Now, for any $x = (x_1, \dots, x_n)$, we may define $d_x(x, y) = \sum_{x_i \neq y_i} x_i$
 (d_H corresponds to $\forall i, x_i = 1$). Using the same method we
 may obtain $\mathbb{P}(x \in A) \mathbb{P}(d_x(x, A) \geq t) \leq e^{-\frac{t^2}{2 \sum x_i^2}}$.

Talagrand's ineq. (no pf.) Ω, x, A like before. Define
 $\rho(x, A) = \sup_{|A|=1} \min_{y \in A} d_x(x, y)$. In particular, $\rho(x, A) \geq d_x(x, A)$. Then
 $\forall A, \forall t \geq 0, \mathbb{P}(x \in A) \mathbb{P}(\rho(x, A) \geq t) \leq e^{-t^2/4}$.

Lemma (Talagrand) $\int_{\Omega} \exp\left(\frac{1}{4} \rho^2(x, A)\right) d\mu \leq \frac{1}{\mathbb{P}(x \in A)}$. pf by induction
 on n using an alternative definition of ρ as a distance
 from the convex hull of a set in the product space. Given the
 lemma, the bound is just Chernoff bounding.